

# Abstract Data Types

EECS 214, Fall 2017

# What is an ADT?

An ADT defines:

- A set of (abstract) values
- A set of (abstract) operations on those values

# What is an ADT?

An ADT defines:

- A set of (abstract) values
- A set of (abstract) operations on those values

An ADT omits:

- How the values are concretely represented
- How the operations work

# ADT: Stack

Looks like: |3 4 5⟩

# ADT: Stack

Looks like: |3 4 5⟩

Signature:

- *push*(Stack, Element): Void
- *pop*(Stack): Element
- *isEmpty*(Stack): Bool

## ADT: Queue (FIFO)

Looks like:  $\langle 3\ 4\ 5 \langle$

## ADT: Queue (FIFO)

Looks like: <3 4 5<

Signature:

- *enqueue*(Queue, Element): Void
- *dequeue*(Queue): Element
- *isEmpty*(Queue): Bool

# Stack versus Queue

Stack signature:

- *push*(Stack, Element): Void
- *pop*(Stack): Element
- *isEmpty*(Stack): Bool

Queue signature:

- *enqueue*(Queue, Element): Void
- *dequeue*(Queue): Element
- *isEmpty*(Queue): Bool



## Adding laws

$$\{p\} \quad f(x) \Rightarrow y \quad \{q\}$$

means that if precondition  $p$  is true when we apply  $f$  to  $x$  then we will get  $y$  as a result, and postcondition  $q$  will be true afterward.

## Adding laws

$$\{p\} \quad f(x) \Rightarrow y \quad \{q\}$$

means that if precondition  $p$  is true when we apply  $f$  to  $x$  then we will get  $y$  as a result, and postcondition  $q$  will be true afterward.

Examples:

$$\{a = [2, 4, 6, 8]\} \quad a[2] \Rightarrow 6 \quad \{a = [2, 4, 6, 8]\}$$

$$\{a = [2, 4, 6, 8]\} \quad a[2] = 0 \quad \{a = [2, 4, 0, 8]\}$$

# ADT: Stack

Looks like:  $|3\ 4\ 5\rangle$

Signature:

- $push(\text{Stack}, \text{Element}): \text{Void}$
- $pop(\text{Stack}): \text{Element}$
- $isEmpty(\text{Stack}): \text{Bool}$

Laws:

$$isEmpty(|\rangle) \Rightarrow \top$$

$$isEmpty(|e_1 \dots e_k e_{k+1}\rangle) \Rightarrow \perp$$

$$\{s = |e_1 \dots e_k\rangle\} push(s, e) \{s = |e_1 \dots e_k e\rangle\}$$

$$\{s = |e_1 \dots e_k e_{k+1}\rangle\} pop(s) \Rightarrow e_{k+1} \{s = |e_1 \dots e_k\rangle\}$$

## ADT: Queue (FIFO)

Looks like:  $\langle 3\ 4\ 5 \langle$

Signature:

- $enqueue(\text{Queue}, \text{Element}): \text{Void}$
- $dequeue(\text{Queue}): \text{Element}$
- $isEmpty(\text{Queue}): \text{Bool}$

Laws:

$$isEmpty(\langle \rangle) \Rightarrow \top$$

$$isEmpty(\langle e_1 \dots e_k e_{k+1} \rangle) \Rightarrow \perp$$

$$\{q = \langle e_1 \dots e_k \rangle\} enqueue(q, e) \{q = \langle e_1 \dots e_k e \rangle\}$$

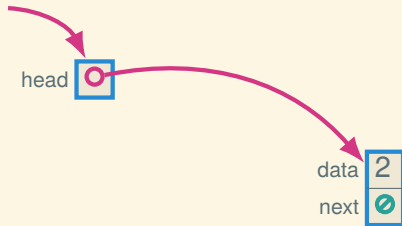
$$\{q = \langle e_1 e_2 \dots e_k \rangle\} dequeue(q) \Rightarrow e_1 \{q = \langle e_2 \dots e_k \rangle\}$$

# Stack implementation: linked list



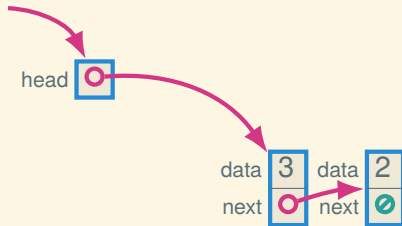
```
let s = new_stack()
```

# Stack implementation: linked list



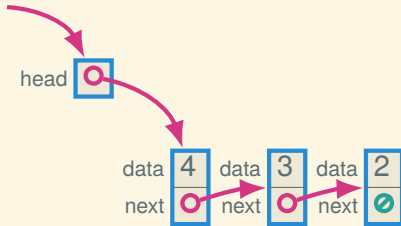
```
let s = new_stack()  
push(s, 2)
```

# Stack implementation: linked list



```
let s = new_stack()  
push(s, 2)  
push(s, 3)
```

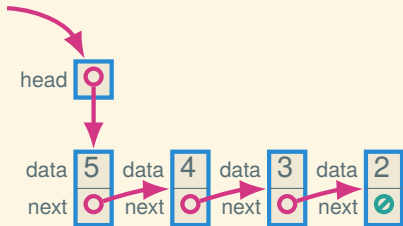
# Stack implementation: linked list



```
let s = new_stack()  
push(s, 2)  
push(s, 3)  
push(s, 4)
```

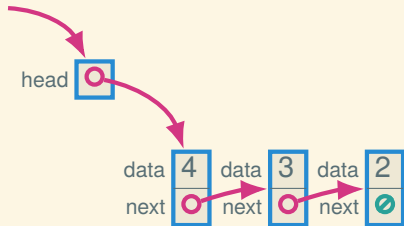


# Stack implementation: linked list



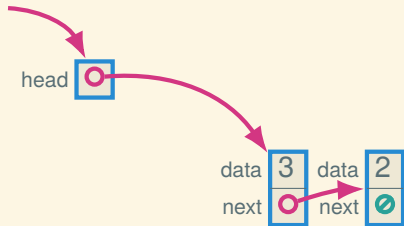
```
let s = new_stack()  
push(s, 2)  
push(s, 3)  
push(s, 4)  
push(s, 5)
```

# Stack implementation: linked list



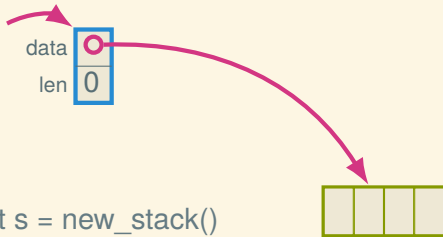
```
let s = new_stack()  
push(s, 2)  
push(s, 3)  
push(s, 4)  
push(s, 5)  
pop(s)
```

# Stack implementation: linked list

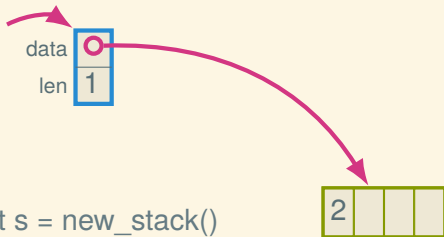


```
let s = new_stack()  
push(s, 2)  
push(s, 3)  
push(s, 4)  
push(s, 5)  
pop(s)  
pop(s)
```

# Stack implementation: array

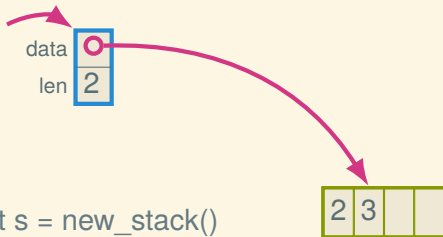


# Stack implementation: array



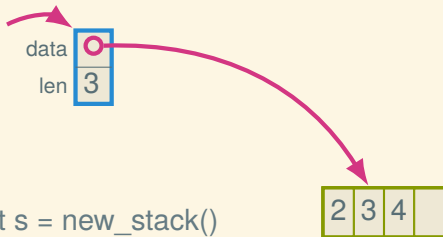
```
let s = new_stack()  
push(s, 2)
```

# Stack implementation: array



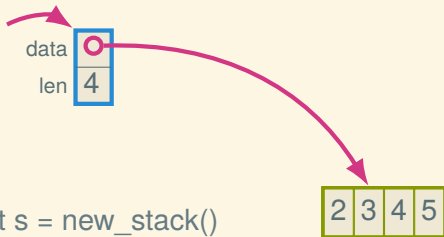
```
let s = new_stack()  
push(s, 2)  
push(s, 3)
```

# Stack implementation: array



```
let s = new_stack()  
push(s, 2)  
push(s, 3)  
push(s, 4)
```

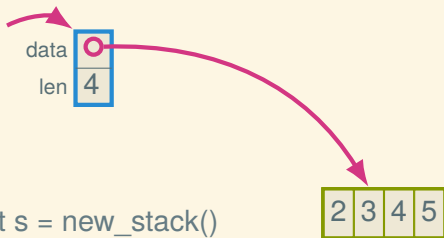
# Stack implementation: array



```
let s = new_stack()  
push(s, 2)  
push(s, 3)  
push(s, 4)  
push(s, 5)
```



# Stack implementation: array



```
let s = new_stack()  
push(s, 2)  
push(s, 3)  
push(s, 4)  
push(s, 5)  
push(s, 6)
```

## ADT: Stack

Looks like:  $|3\ 4\ 5\rangle$

Signature:

- $push(\text{Stack}, \text{Element}): \text{Void} \text{ — } \mathcal{O}(1)$
- $pop(\text{Stack}): \text{Element} \text{ — } \mathcal{O}(1)$
- $isEmpty(\text{Stack}): \text{Bool} \text{ — } \mathcal{O}(1)$

Laws:

$$isEmpty(|\rangle) \Rightarrow \top$$

$$isEmpty(|e_1 \dots e_k e_{k+1}\rangle) \Rightarrow \perp$$

$$\{s = |e_1 \dots e_k\rangle\} push(s, e) \{s = |e_1 \dots e_k e\rangle\}$$

$$\{s = |e_1 \dots e_k e_{k+1}\rangle\} pop(s) \Rightarrow e_{k+1} \{s = |e_1 \dots e_k\rangle\}$$

## Trade-offs: linked list stack versus array stack

- Linked list stack only fills up when memory fills up, whereas array stack has a fixed size (or must reallocate)
- Array stack has better constant factors: cache locality and no (or rare) allocation
- Array stack space usage is tighter; linked list is smoother

## ADT: Queue (FIFO)

Looks like:  $\langle 3\ 4\ 5 \langle$

Signature:

- $enqueue(\text{Queue}, \text{Element}): \text{Void} \text{ — } \mathcal{O}(1)$
- $dequeue(\text{Queue}): \text{Element} \text{ — } \mathcal{O}(1)$
- $isEmpty(\text{Queue}): \text{Bool} \text{ — } \mathcal{O}(1)$

Laws:

$$isEmpty(\langle \rangle) \Rightarrow \top$$

$$isEmpty(\langle e_1 \dots e_k e_{k+1} \rangle) \Rightarrow \perp$$

$$\{q = \langle e_1 \dots e_k \rangle\} enqueue(q, e) \{q = \langle e_1 \dots e_k e \rangle\}$$

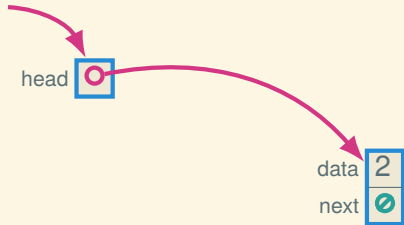
$$\{q = \langle e_1 e_2 \dots e_k \rangle\} dequeue(q) \Rightarrow e_1 \{q = \langle e_2 \dots e_k \rangle\}$$

## Queue implementation: linked list?



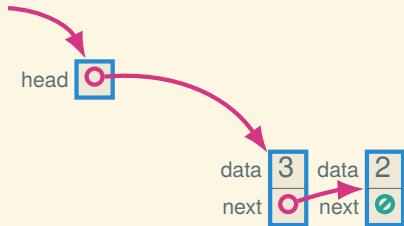
```
let q = new_queue()
```

## Queue implementation: linked list?



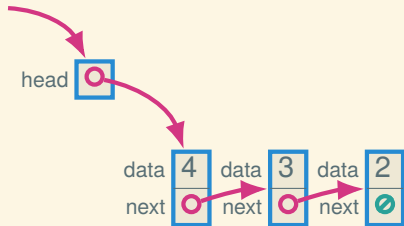
```
let q = new_queue()  
enqueue(q, 2)
```

## Queue implementation: linked list?



```
let q = new_queue()  
enqueue(q, 2)  
enqueue(q, 3)
```

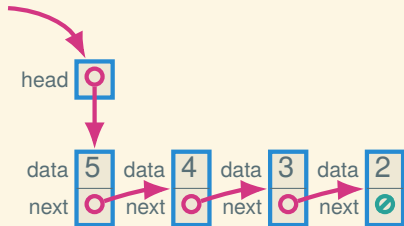
## Queue implementation: linked list?



```
let q = new_queue()  
enqueue(q, 2)  
enqueue(q, 3)  
enqueue(q, 4)
```

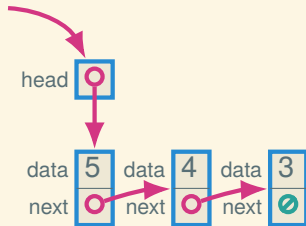


## Queue implementation: linked list?



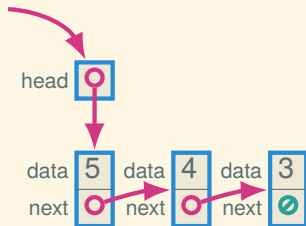
```
let q = new_queue()  
enqueue(q, 2)  
enqueue(q, 3)  
enqueue(q, 4)  
enqueue(q, 5)
```

## Queue implementation: linked list?



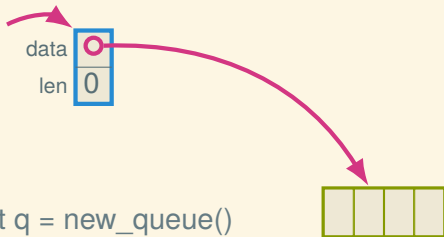
```
let q = new_queue()  
enqueue(q, 2)  
enqueue(q, 3)  
enqueue(q, 4)  
enqueue(q, 5)  
dequeue(q)
```

## Queue implementation: linked list?



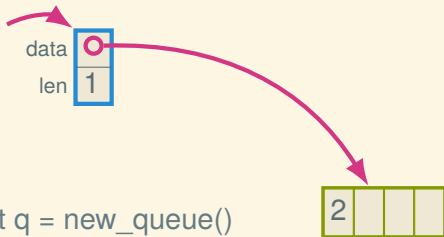
```
let q = new_queue()  
enqueue(q, 2)  
enqueue(q, 3)  
enqueue(q, 4)  
enqueue(q, 5)  
dequeue(q) —  $O(n)$ ?
```

# Queue implementation: array?



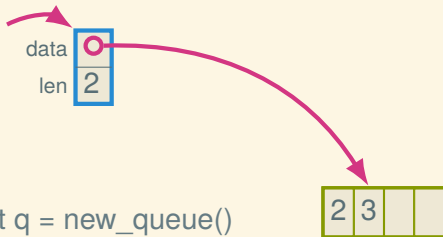
let q = new\_queue()

# Queue implementation: array?



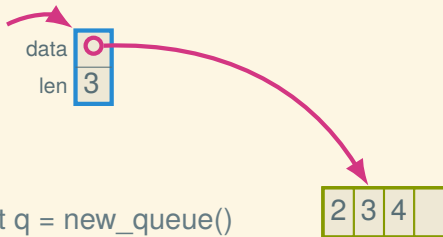
```
let q = new_queue()  
enqueue(q, 2)
```

# Queue implementation: array?



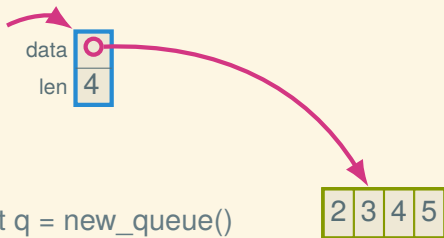
```
let q = new_queue()  
enqueue(q, 2)  
enqueue(q, 3)
```

## Queue implementation: array?



```
let q = new_queue()  
enqueue(q, 2)  
enqueue(q, 3)  
enqueue(q, 4)
```

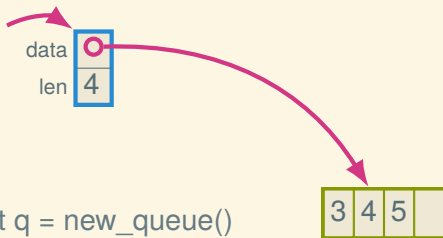
## Queue implementation: array?



```
let q = new_queue()  
enqueue(q, 2)  
enqueue(q, 3)  
enqueue(q, 4)  
enqueue(q, 5)
```

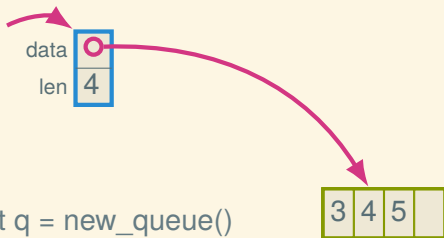


## Queue implementation: array?



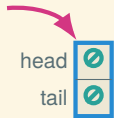
```
let q = new_queue()  
enqueue(q, 2)  
enqueue(q, 3)  
enqueue(q, 4)  
enqueue(q, 5)  
s.dequeue()
```

## Queue implementation: array?



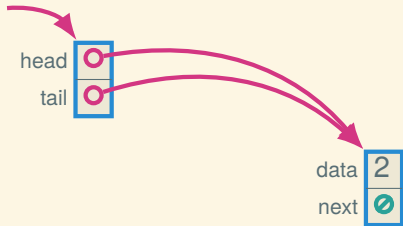
```
let q = new_queue()
enqueue(q, 2)
enqueue(q, 3)
enqueue(q, 4)
enqueue(q, 5)
s.dequeue() —  $\mathcal{O}(n)$ ???
```

## Queue impl.: linked list with tail pointer



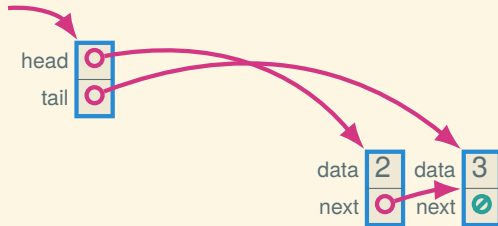
```
let q = new_queue()
```

## Queue impl.: linked list with tail pointer



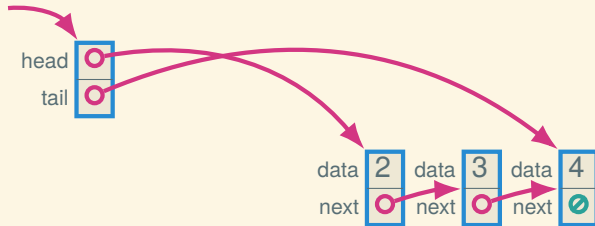
```
let q = new_queue()  
enqueue(q, 2)
```

## Queue impl.: linked list with tail pointer



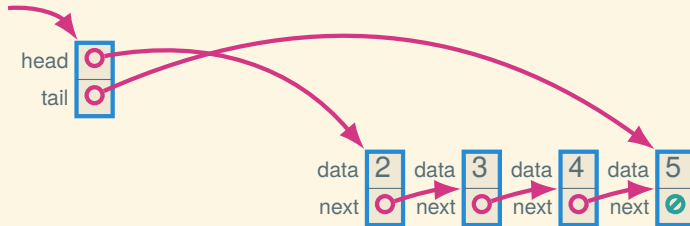
```
let q = new_queue()  
enqueue(q, 2)  
enqueue(q, 3)
```

# Queue impl.: linked list with tail pointer



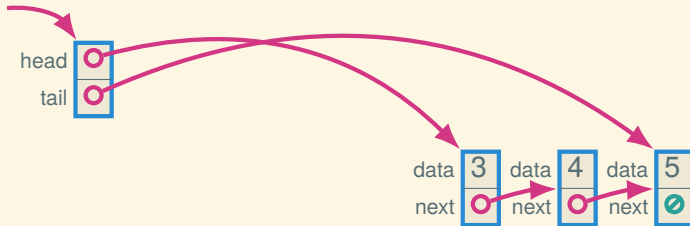
```
let q = new_queue()  
enqueue(q, 2)  
enqueue(q, 3)  
enqueue(q, 4)
```

## Queue impl.: linked list with tail pointer



```
let q = new_queue()  
enqueue(q, 2)  
enqueue(q, 3)  
enqueue(q, 4)  
enqueue(q, 5)
```

## Queue impl.: linked list with tail pointer



```
let q = new_queue()
enqueue(q, 2)
enqueue(q, 3)
enqueue(q, 4)
enqueue(q, 5)
dequeue(q)
```

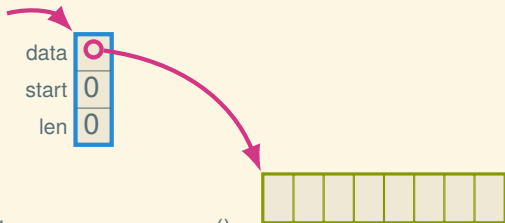


## Queue impl.: linked list with tail pointer



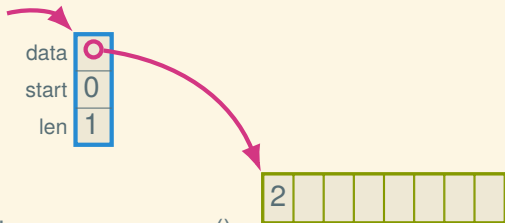
```
let q = new_queue()
enqueue(q, 2)
enqueue(q, 3)
enqueue(q, 4)
enqueue(q, 5)
dequeue(q)
dequeue(q)
```

# Queue implementation: ring buffer



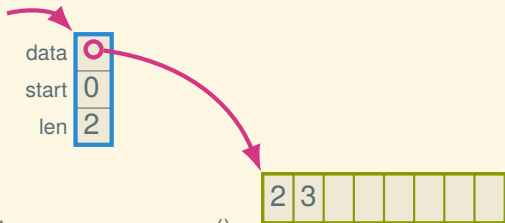
```
let q = new_queue()
```

# Queue implementation: ring buffer



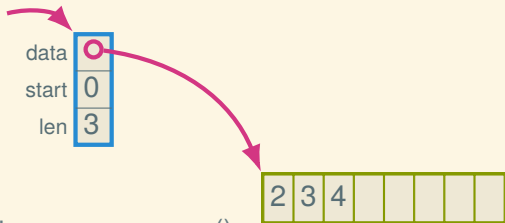
```
let q = new_queue()  
enqueue(q, 2)
```

# Queue implementation: ring buffer



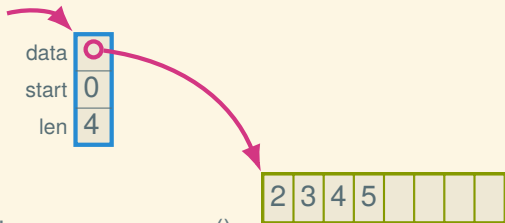
```
let q = new_queue()  
enqueue(q, 2)  
enqueue(q, 3)
```

# Queue implementation: ring buffer



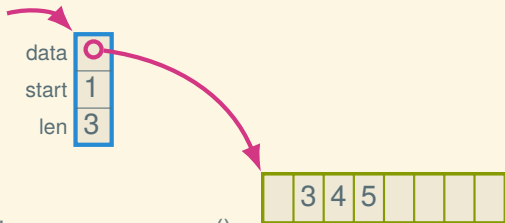
```
let q = new_queue()  
enqueue(q, 2)  
enqueue(q, 3)  
enqueue(q, 4)
```

# Queue implementation: ring buffer



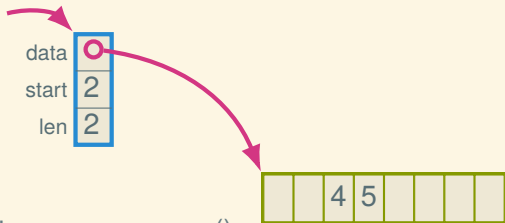
```
let q = new_queue()  
enqueue(q, 2)  
enqueue(q, 3)  
enqueue(q, 4)  
enqueue(q, 5)
```

# Queue implementation: ring buffer



```
let q = new_queue()  
enqueue(q, 2)  
enqueue(q, 3)  
enqueue(q, 4)  
enqueue(q, 5)  
dequeue(q)
```

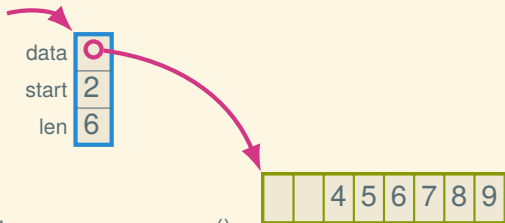
# Queue implementation: ring buffer



```
let q = new_queue()  
enqueue(q, 2)  
enqueue(q, 3)  
enqueue(q, 4)  
enqueue(q, 5)  
dequeue(q)  
dequeue(q)
```

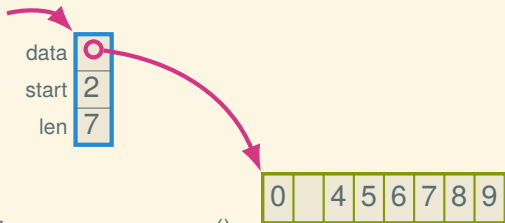


# Queue implementation: ring buffer



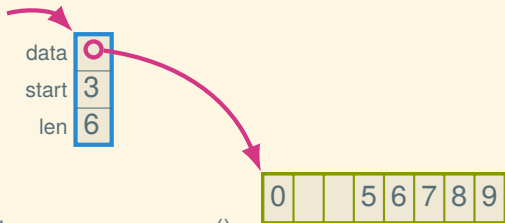
```
let q = new_queue()  
enqueue(q, 2)  
enqueue(q, 3)  
enqueue(q, 4)  
enqueue(q, 5)  
dequeue(q)  
dequeue(q)  
⋮
```

# Queue implementation: ring buffer



```
let q = new_queue()  
enqueue(q, 2)  
enqueue(q, 3)  
enqueue(q, 4)  
enqueue(q, 5)  
dequeue(q)  
dequeue(q)  
:  
enqueue(q, 0)
```

# Queue implementation: ring buffer



```
let q = new_queue()  
enqueue(q, 2)  
enqueue(q, 3)  
enqueue(q, 4)  
enqueue(q, 5)  
dequeue(q)  
dequeue(q)  
:  
enqueue(q, 0)  
dequeue(q)
```

## Trade-offs: linked list queue versus ring buffer

Basically the same as for the stack implementations:

- Ring buffer has better constant factors and uses less space (potentially)
- Linked list doesn't fill up

## Ring buffer in DSSL2

# Representation

```
# A QueueOf[X] is
#   queue(VectorOf[X or False], Natural, Natural)
# Interpretation:
# - `data` contains the elements of the queue,
# - `start` is the index of the first element, and
# - `size` is the number of elements.
defstruct queue(data, start, size)
```

# Representation

```
# A QueueOf[X] is  
#   queue(VectorOf[X or False], Natural, Natural)  
# Interpretation:  
# - `data` contains the elements of the queue,  
# - `start` is the index of the first element, and  
# - `size` is the number of elements.  
defstruct queue(data, start, size)  
  
let F = False  
let QUEUE0 = queue([F; 8], 0, 0)  
let QUEUE1 = queue([3, F, F, F, F, F], 0, 1)  
let QUEUE2 = queue([3, 4, F, F, F, F], 0, 2)  
let QUEUE3 = queue([F, F, 5, 6, 7, F], 2, 3)  
let QUEUE4 = queue([9, 10, F, F, 7, 8], 4, 4)
```

## Creating a new queue

```
# new_queue : Natural -> QueueOf[X]
def new_queue(capacity):
    queue([ False; capacity ], 0, 0)
```



## Finding out the size and capacity

```
# queue_size : QueueOf[X] -> Natural  
def queue_size(q): q.size
```

```
# queue_capacity : QueueOf[X] -> Natural  
def queue_capacity(q): len(q.data)
```

## Finding out the size and capacity

```
# queue_size : QueueOf[X] -> Natural  
def queue_size(q): q.size
```

```
# queue_capacity : QueueOf[X] -> Natural  
def queue_capacity(q): len(q.data)
```

```
# queue_empty? : QueueOf[X] -> Bool  
def queue_empty?(q):  
    queue_size(q) == 0
```

```
# queue_full? : QueueOf[X] -> Bool  
def queue_full?(q):  
    queue_size(q) == queue_capacity(q)
```

# Enqueueing

```
# enqueue! : QueueOf[X] X -> Void
def enqueue!(q, element):
  if queue_full?(q):
    error('enqueue!: queue is full')
  let cap = queue_capacity(q)
  q.data[(q.size + q.start) % cap] = element
  q.size = q.size + 1
```

## Dequeuing

```
# dequeue! : QueueOf[X] -> X
def dequeue!(q):
  if queue_empty?(q):
    error('dequeue!: queue is empty')
  let result = q.data[q.start]
  q.data[q.start] = False
  q.size = q.size - 1
  q.start = (q.start + 1) % queue_capacity(q)
  result
```

Next time: BSTs and the Dictionary ADT