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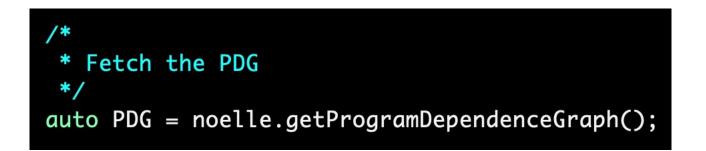
Outline

• Program Dependence Graph at the instruction granularity

• SCCDAG

• Semantics of dependences

PDG* is provided by NOELLE



This PDG is at the instruction granularity

- A dependence is either
 - Between two instructions or
 - Between an instruction and a function parameter

[*] Jeanne Ferrante, Karl J. Ottenstein, Joe D. Warren.

The program dependence graph and its use in optimization. ACM Transactions on Programming Languages and System 1987

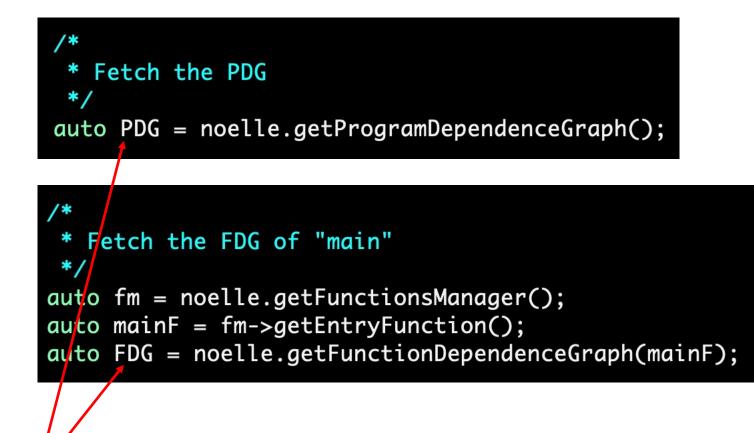
NOELLE's PDG at the instruction granularity

- Dependences are clustered by function
- Dependences between instructions in two functions:

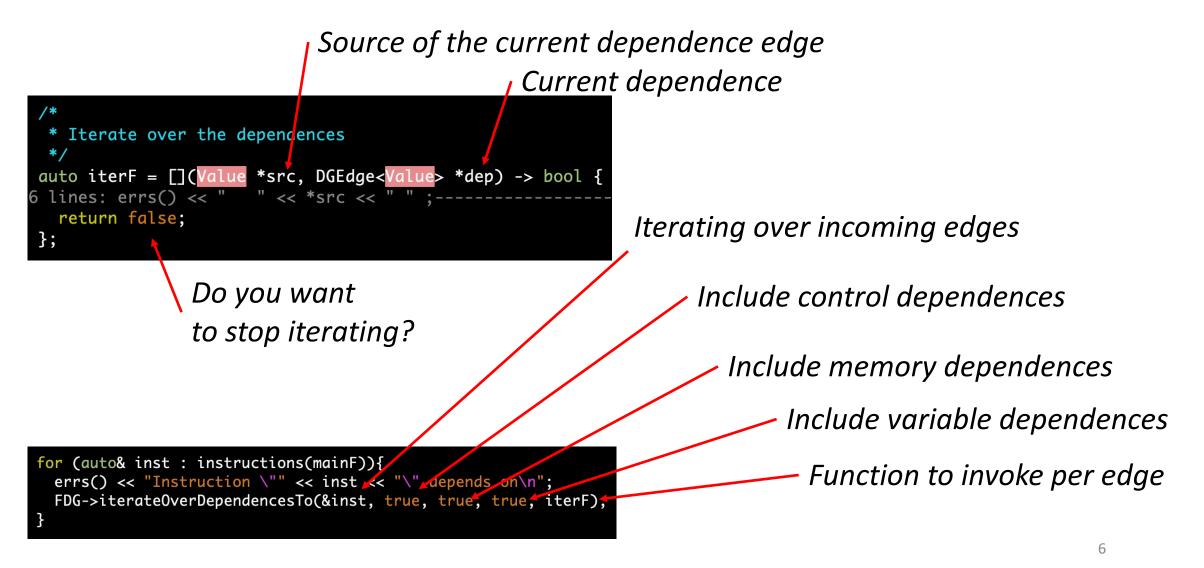
```
declare void @f2 (int8 *%0){
```

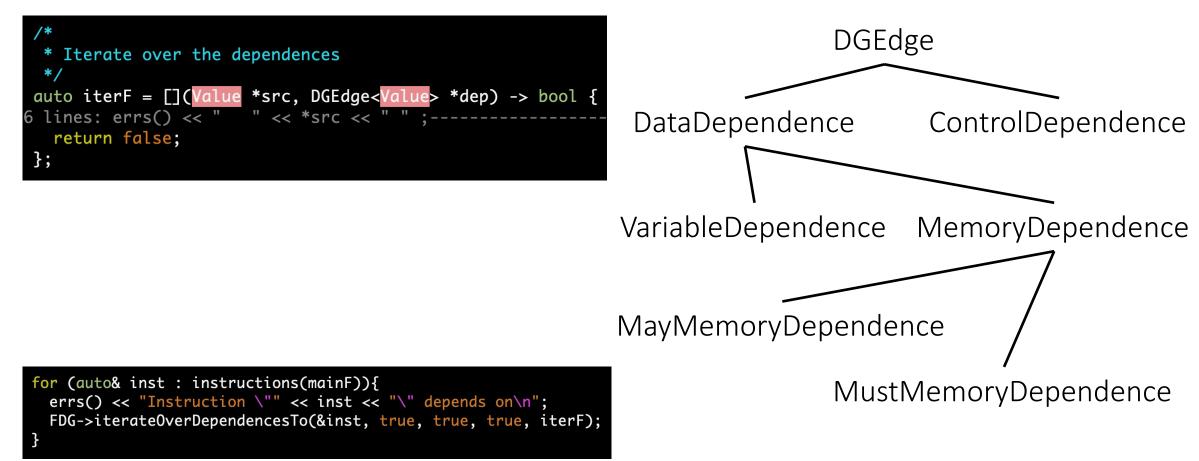
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NOELLE's Function Dependence Graph (FDG)



Different instances of the same C++ class (PDG)

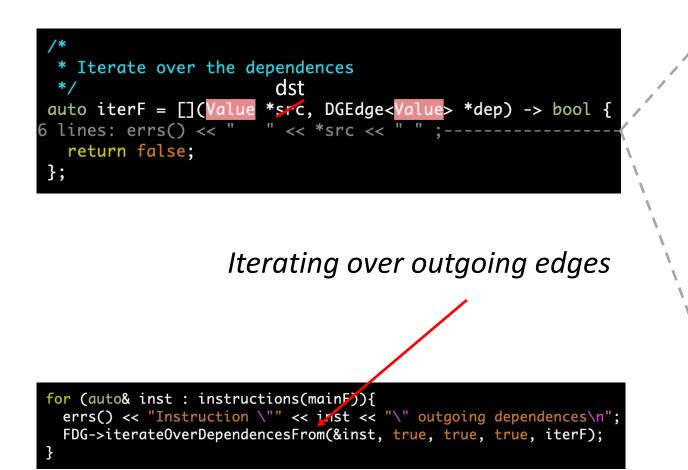




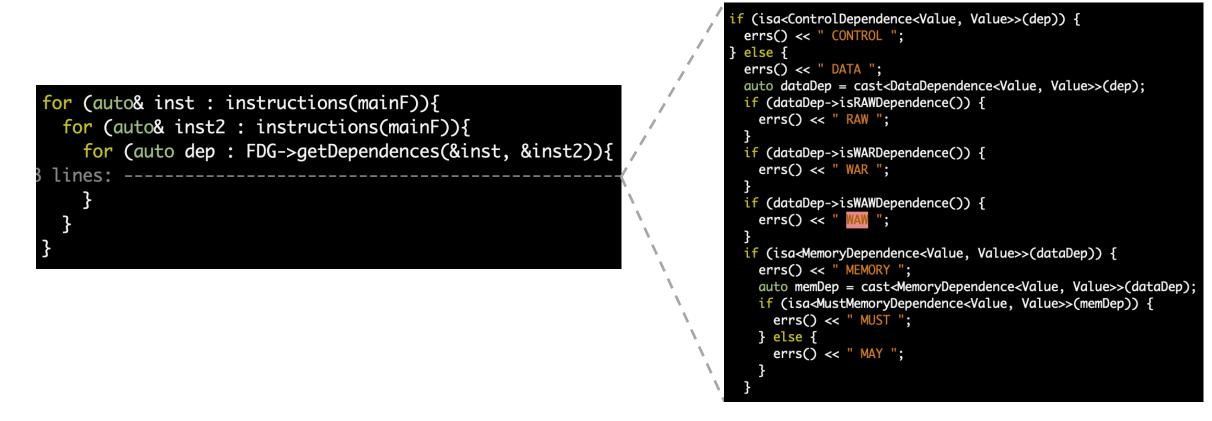


for (auto& inst : instructions(mainF)){
errs() << "Instruction \"" << inst << "\" depends on\n";
FDG->iterateOverDependencesTo(&inst, true, true, true, iterF);









This code can be found here: noelle/examples/passes/pdg

NOELLE provides SCCDAG

- NOELLE provides:
 - Program Dependence Graph (PDG)
 - Function Dependence Graph (FDG)

- Different scope (i.e., code region they target)
- Dependences between dynamic instances of X are not included in XDG but are included in YDG where Y includes X
- Loop Dependence Graph (LDG) (see NOELLE_loops slides/talk)
- All dependence graphs are instances of the same class arcana::noelle::PDG
- Because of importance of loops, NOELLE provides a rich class for them called arcana::noelle::LoopContent
- LoopContent includes:
 - LDG
 - SCCDAG
 - And much more (see NOELLE_loops slides/talk)

Memory alias analysis: the problem (from 323)

• We want to

- Execute *j* in parallel with *i* (extracting parallelism)
- Move *j* before *i* (code scheduling)
- Does *j* depend on *i* ?

- Do p and q point to the same memory location?
 - Does q alias p?



Memory alias analyses included in NOELLE

- NOELLE relies on ~40 memory alias analyses to compute its PDG
- Most analyses are included in the following 3 frameworks:
 - SCAF: <u>https://github.com/PrincetonUniversity/SCAF</u>
 - SVF: <u>https://github.com/SVF-tools/SVF</u>
 - LLVM: <u>http://llvm.org</u>
- NOELLE includes an extra alias analysis as well to capture corner cases that alias analyses above do not
 - We see alias analysis to be used by NOELLE, rather than for NOELLE to provide
 - Hence, when another alias infrastructure will capture them, this NOELLE's AA will be removed

Outline

• Program Dependence Graph at the instruction granularity

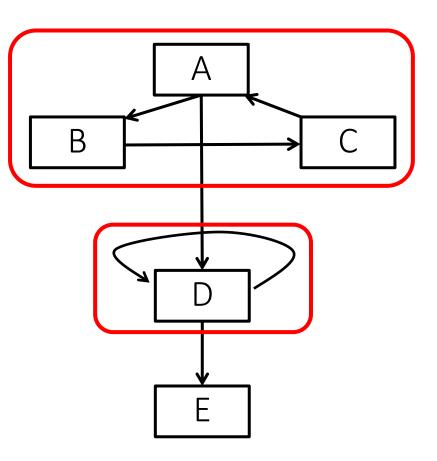


• Semantics of dependences

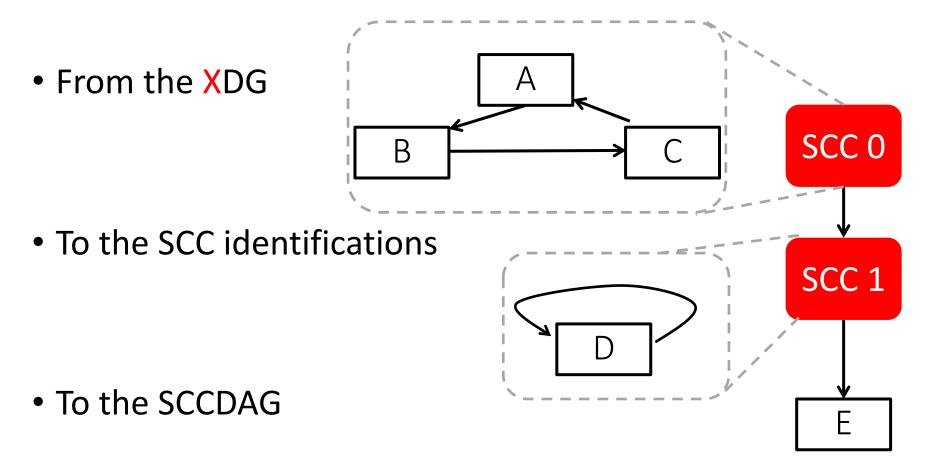
NOELLE's Hierarchical SCCDAG

• From the XDG

• To the SCC identifications



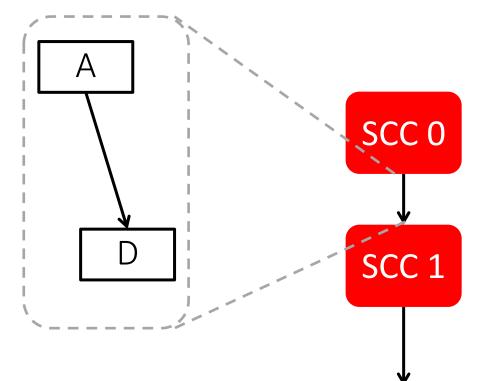
NOELLE's Hierarchical SCCDAG



NOELLE's Hierarchical SCCDAG

• From the XDG

• To the SCC identifications



Ε

• To the SCCDAG



Outline

• Program Dependence Graph at the instruction granularity



• Semantics of dependences

Dependences

- Control dependences
- Data dependences
 - Variable
 - Memory

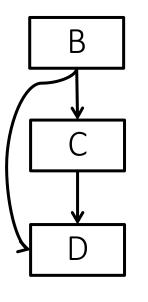
Dependences

- Control dependences
- Data dependences
 - Variable
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 - May
 - Must

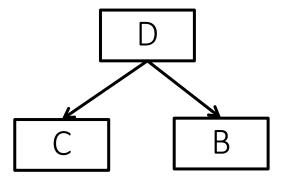
Post-Dominators

Assumption: Single exit node in CFG

Definition: Node *d* post-dominates node *n* in a graph if every path from *n* to the exit node goes through *d*



CFG

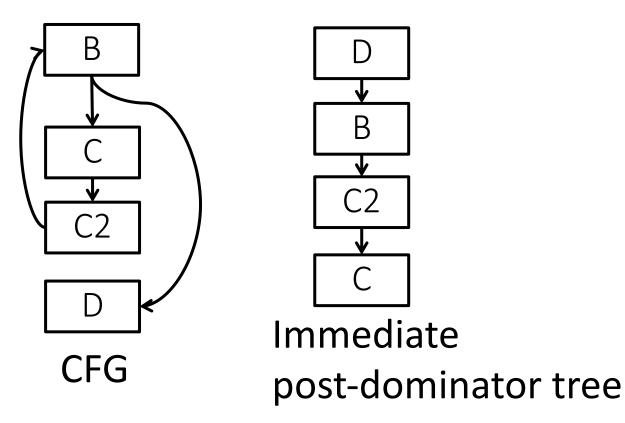


Immediate post-dominator tree

Control dependences

A node Y control-depends on another node X if and only if

- 1. There is a path from X to Y such that every node in that path other than X is post-dominated by Y
- 2. X is not **strictly** post-dominated by Y



Dependences

- Control dependences
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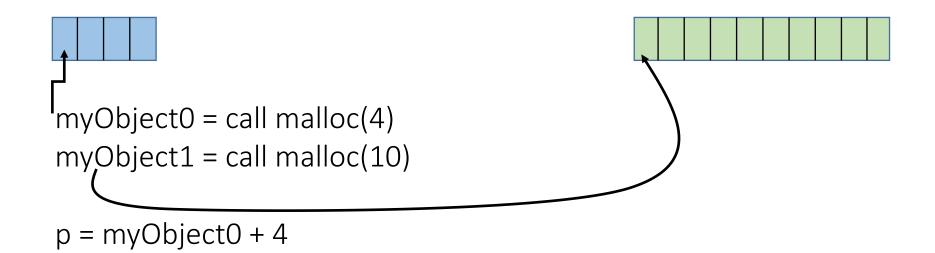
Data dependences

- A variable dependence is a def-use chain in LLVM
- A memory dependence from instruction i1 to instruction i2 exists iff *:
 - the footprint of operation i1 may-alias the footprint of i2 (alias);
 - at least one of the two instructions writes to memory (update);
 - there is a feasible path of execution P from i1 to i2 (feasible-path) such that no operation in P overwrites the common memory footprint (no-kill).

Footprint refers to the memory locations accessed (read or written) by an instruction.

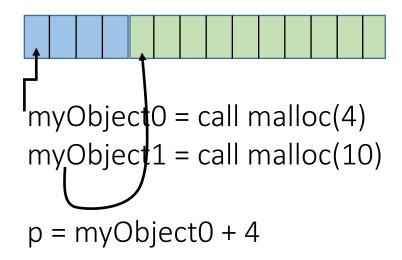
[*] Sotiris Apostolakis , Ziyang Xu , Zujun Tan , Greg Chan, Simone Campanoni, and David I. August SCAF: A Speculation-Aware Collaborative Dependence Analysis Framework. PLDI 2020.

The (LLVM) memory model



Can p alias myObject1?

The (LLVM) memory model



Can p alias myObject1?

Always have faith in your ability

Success will come your way eventually

Best of luck!